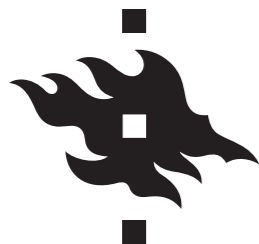


Boolean matrix factorization *meets* consecutive ones property

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Boolean matrix factorisation

- Given a matrix **A**, find matrices **B** and **C** s.t. $\mathbf{A} \approx \mathbf{B} \circ \mathbf{C}^T$
 - SVD, NMF, PCA, ICA, ...
- In **Boolean matrix factorisation** (BMF)
 1. the input matrix **A** is binary
 2. the factor matrices **B** and **C** are binary
 3. the algebra is Boolean ($1 + 1 = 1$)
- Sparse factors, good interpretability, sometimes you just need binary factors

Example of BMF

$$\begin{array}{c} 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{array} \begin{array}{c} 1 \\ 2 \\ 3 \\ 4 \end{array} \begin{pmatrix} 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \end{pmatrix} = \begin{array}{c} 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{array} \begin{array}{c} a \\ b \\ c \end{array} \begin{pmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix} \circ \begin{array}{c} d \\ e \\ f \end{array} \begin{array}{c} 1 \\ 2 \\ 3 \\ 4 \end{array} \begin{pmatrix} 1 & 2 & 3 & 4 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 \end{pmatrix}$$

Example of BMF

4th column = $a + b$

$$\begin{array}{c}
 1 \\
 2 \\
 3 \\
 4 \\
 5 \\
 6
 \end{array}
 \begin{array}{cccc}
 1 & 2 & 3 & 4 \\
 \left(\begin{array}{cccc}
 1 & 0 & 1 & 1 \\
 1 & 1 & 1 & 1 \\
 0 & 1 & 1 & 1 \\
 0 & 1 & 1 & 1 \\
 0 & 1 & 1 & 0 \\
 1 & 0 & 1 & 1
 \end{array} \right)
 \end{array}
 =
 \begin{array}{c}
 1 \\
 2 \\
 3 \\
 4 \\
 5 \\
 6
 \end{array}
 \begin{array}{ccc}
 a & b & c \\
 \left(\begin{array}{ccc}
 1 & 0 & 0 \\
 1 & 1 & 0 \\
 0 & 1 & 0 \\
 0 & 1 & 0 \\
 0 & 0 & 1 \\
 1 & 0 & 0
 \end{array} \right)
 \end{array}
 \circ
 \begin{array}{c}
 d \\
 e \\
 f
 \end{array}
 \begin{array}{cccc}
 1 & 2 & 3 & 4 \\
 \left(\begin{array}{cccc}
 1 & 0 & 1 & 1 \\
 0 & 1 & 1 & 1 \\
 0 & 1 & 1 & 0
 \end{array} \right)
 \end{array}$$

Example of BMF

4th column = $a + b$

$$\begin{array}{cccc}
 & 1 & 2 & 3 & 4 \\
 1 & 1 & 0 & 1 & 1 \\
 2 & 1 & 1 & 1 & 1 \\
 3 & 0 & 1 & 1 & 1 \\
 4 & 0 & 1 & 1 & 1 \\
 5 & 0 & 1 & 1 & 0 \\
 6 & 1 & 0 & 1 & 1
 \end{array}
 =
 \begin{array}{ccc}
 & a & b & c \\
 1 & 1 & 0 & 0 \\
 2 & 1 & 1 & 0 \\
 3 & 0 & 1 & 0 \\
 4 & 0 & 1 & 0 \\
 5 & 0 & 0 & 1 \\
 6 & 1 & 0 & 0
 \end{array}
 \circ
 \begin{array}{cccc}
 & 1 & 2 & 3 & 4 \\
 d & 1 & 0 & 1 & 1 \\
 e & 0 & 1 & 1 & 1 \\
 f & 0 & 1 & 1 & 0
 \end{array}$$

$(2, 4) = 1 = 1 + 1 = (2, a) + (2, b)$

Consecutive ones property

- A binary vector \mathbf{x} has the **consecutive ones property** (C1P) if all of its 1s are consecutive
 - e.g. $\mathbf{x} = (0, 0, 1, 1, 1, 1, 1, 0)$
- A binary matrix \mathbf{X} has the C1P if its rows can be permuted so that all of its columns have C1P simultaneously
- A BMF $\mathbf{A} \approx \mathbf{B} \circ \mathbf{C}^T$ has the C1P if both \mathbf{B} and \mathbf{C} have the C1P

Cyclic vectors and matrices

- A binary vector \mathbf{x} is **cyclic** if it's complement has the consecutive ones property
 - e.g. $\mathbf{x} = (1, 1, 0, 0, 0, 0, 0, 1)$
- A binary matrix is **cyclic** if its columns have either the C1P or the cyclic ones property
- A BMF $\mathbf{A} \approx \mathbf{B} \circ \mathbf{C}^T$ is **cyclic** if both \mathbf{B} and \mathbf{C} are cyclic

Example

$$\begin{array}{c} 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{array} \begin{array}{cccc} 1 & 2 & 3 & 4 \\ \left(\begin{array}{cccc} 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \end{array} \right) \end{array} = \begin{array}{c} 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{array} \begin{array}{ccc} a & b & c \\ \left(\begin{array}{ccc} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{array} \right) \end{array} \circ \begin{array}{c} d \\ e \\ f \end{array} \begin{array}{cccc} 1 & 2 & 3 & 4 \\ \left(\begin{array}{cccc} 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 \end{array} \right) \end{array}$$

Example

$$\begin{array}{c} 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{array} \begin{array}{cccc} 1 & 2 & 3 & 4 \\ \left(\begin{array}{cccc} 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \end{array} \right) \end{array} = \begin{array}{c} \text{cyclic} \\ 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{array} \begin{array}{ccc} a & b & c \\ \left(\begin{array}{ccc} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{array} \right) \end{array} \circ \begin{array}{c} d \\ e \\ f \end{array} \begin{array}{cccc} 1 & 2 & 3 & 4 \\ \left(\begin{array}{cccc} 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 \end{array} \right) \end{array}$$

Example

$$\begin{array}{c} 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{array} \begin{array}{cccc} 1 & 2 & 3 & 4 \\ \left(\begin{array}{cccc} 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \end{array} \right) \end{array} = \begin{array}{c} 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{array} \begin{array}{ccc} \text{cyclic CIP} \\ a & b & c \\ \left(\begin{array}{ccc} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{array} \right) \end{array} \circ \begin{array}{c} d \\ e \\ f \end{array} \begin{array}{cccc} 1 & 2 & 3 & 4 \\ \left(\begin{array}{cccc} 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 \end{array} \right) \end{array}$$

Example

$$\begin{array}{c}
 \\
 \\
 \\
 \\
 \\
 \\
 \end{array}
 \begin{array}{cccc}
 1 & 2 & 3 & 4 \\
 \left(\begin{array}{cccc}
 1 & 0 & 1 & 1 \\
 1 & 1 & 1 & 1 \\
 0 & 1 & 1 & 1 \\
 0 & 1 & 1 & 1 \\
 0 & 1 & 1 & 0 \\
 1 & 0 & 1 & 1
 \end{array} \right)
 \end{array}
 =
 \begin{array}{c}
 \\
 \\
 \\
 \\
 \\
 \\
 \end{array}
 \begin{array}{ccc}
 \text{cyclic C1P} & \text{C1P} & \text{C1P} \\
 a & b & c \\
 \left(\begin{array}{ccc}
 1 & 0 & 0 \\
 1 & 1 & 0 \\
 0 & 1 & 0 \\
 0 & 1 & 0 \\
 0 & 0 & 1 \\
 1 & 0 & 0
 \end{array} \right)
 \end{array}
 \circ
 \begin{array}{c}
 d \\
 e \\
 f
 \end{array}
 \begin{array}{cccc}
 1 & 2 & 3 & 4 \\
 \left(\begin{array}{cccc}
 1 & 0 & 1 & 1 \\
 0 & 1 & 1 & 1 \\
 0 & 1 & 1 & 0
 \end{array} \right)
 \end{array}$$

Example

$$\begin{array}{c}
 1 \\
 2 \\
 3 \\
 4 \\
 5 \\
 6
 \end{array}
 \begin{array}{c}
 1 \\
 2 \\
 3 \\
 4
 \end{array}
 \begin{pmatrix}
 1 & 0 & 1 & 1 \\
 1 & 1 & 1 & 1 \\
 0 & 1 & 1 & 1 \\
 0 & 1 & 1 & 1 \\
 0 & 1 & 1 & 0 \\
 1 & 0 & 1 & 1
 \end{pmatrix}
 =
 \begin{array}{c}
 1 \\
 2 \\
 3 \\
 4 \\
 5 \\
 6
 \end{array}
 \begin{array}{c}
 \text{cyclic} \\
 \text{C1P} \\
 \text{C1P}
 \end{array}
 \begin{array}{c}
 a \\
 b \\
 c
 \end{array}
 \begin{pmatrix}
 1 & 0 & 0 \\
 1 & 1 & 0 \\
 0 & 1 & 0 \\
 0 & 1 & 0 \\
 0 & 0 & 1 \\
 1 & 0 & 0
 \end{pmatrix}
 \circ
 \begin{array}{c}
 d \\
 e \\
 f
 \end{array}
 \begin{array}{c}
 1 \\
 2 \\
 3 \\
 4
 \end{array}
 \begin{pmatrix}
 1 & 0 & 1 & 1 \\
 1 & 0 & 1 & 1 \\
 0 & 1 & 1 & 1 \\
 0 & 1 & 1 & 0
 \end{pmatrix}
 \begin{array}{c}
 \text{cyclic} \\
 \text{C1P} \\
 \text{C1P}
 \end{array}$$

Why?

- We aim at finding cyclic or C1P BMF
 - All problems are NP-hard
 - The extra constraints typically make the reconstruction less accurate
 - But they can also make the factors easier to understand
 - E.g. rows are time points
 - ⇒ C1P = something starts, lasts, and ends, but doesn't resurrect
- And we can draw graphs

How?

Typical BMF algorithm:

1. construct a set of candidate vectors

2. repeat

2.1. add a new factor from the
candidates based on what reduces
the error most

3. until k factors are selected

But how to satisfy the constraints?

PQ-trees

- A PQ-tree is a data structure from the '70s that encodes permutations
 - the leaves are the indices
 - the children of P-nodes can be ordered arbitrarily
 - the children of Q-nodes must stay in order (or reverse)

Theorem (Booth & Lueker, '76). There exists a PQ-tree that encodes exactly the row permutations where a matrix has the C1P (or cyclic) property and this tree can be updated in polynomial time.

Using PQ-trees

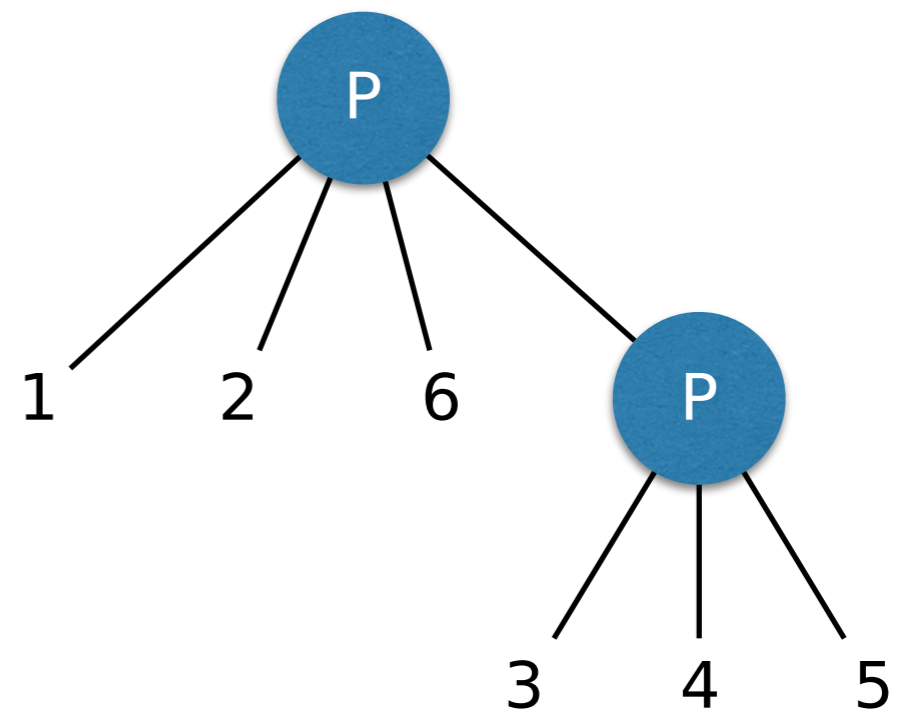
- We keep a PQ-tree for both factor matrices
- Every candidate column introduces new constraints
 - Update the PQ-tree with these constraints
 - ⇒ will show if constraints can't be satisfied
- Select the best of the possible candidates
 - Quality can be calculated quickly
- One factor can be added in $O(\text{nnz}(\mathbf{A}) + m + n)$ time

Example

$$\begin{array}{c} a \\ 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{array} \begin{pmatrix} 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \end{pmatrix}$$

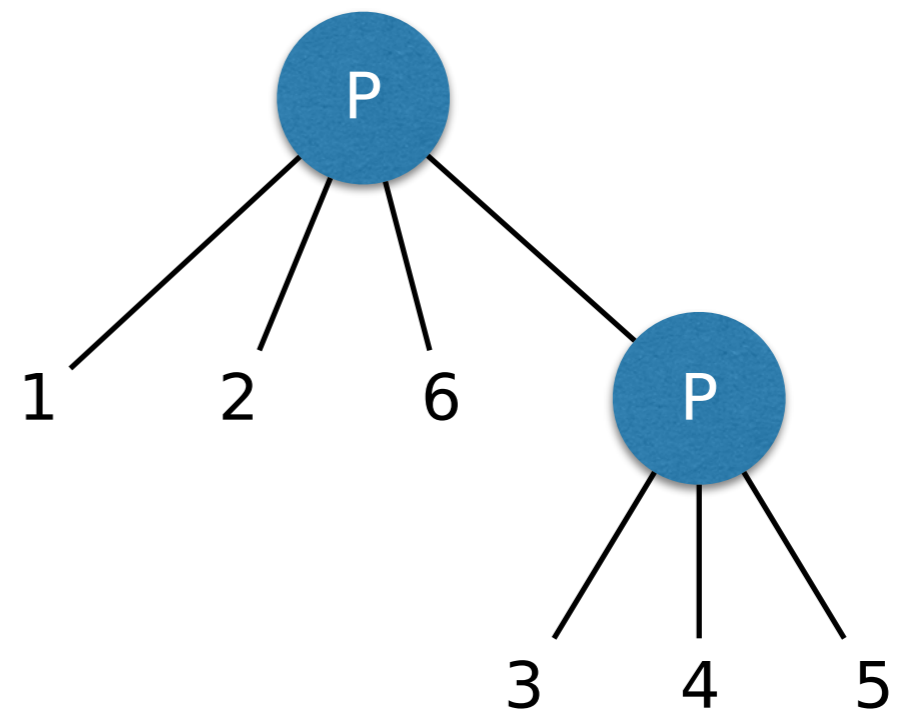
Example

$$\begin{array}{c} 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{array} \begin{array}{c} a \\ \left(\begin{array}{c} 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \end{array} \right) \end{array}$$



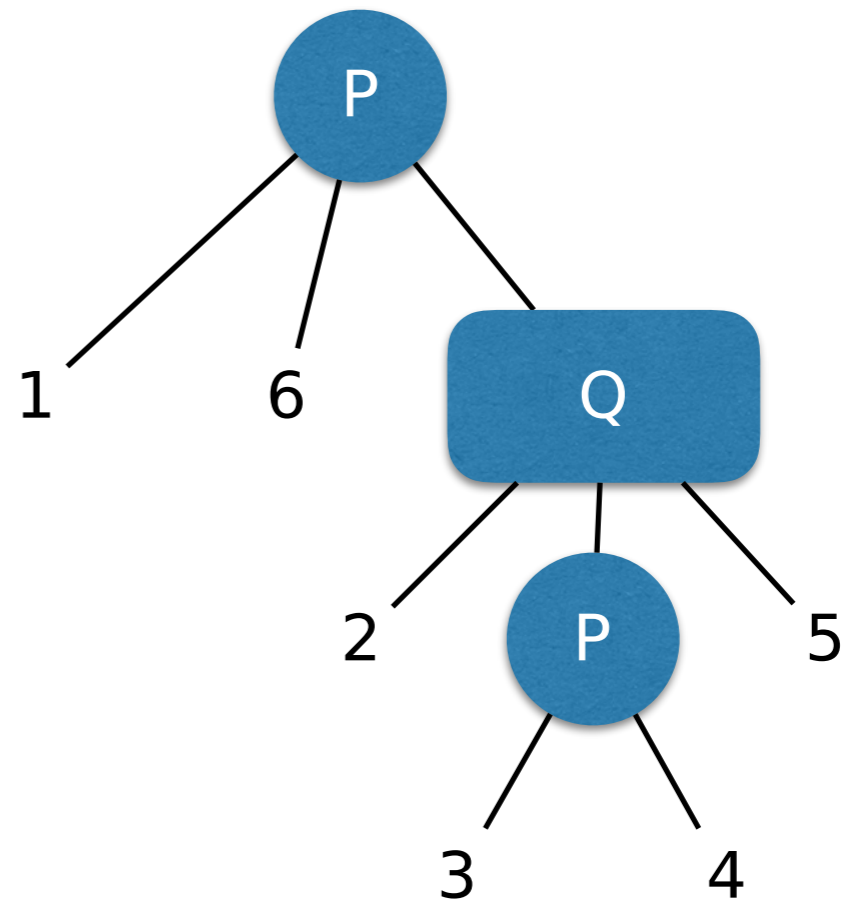
Example

	<i>a</i>	<i>b</i>
1	1	0
2	1	1
3	0	1
4	0	1
5	0	0
6	1	0

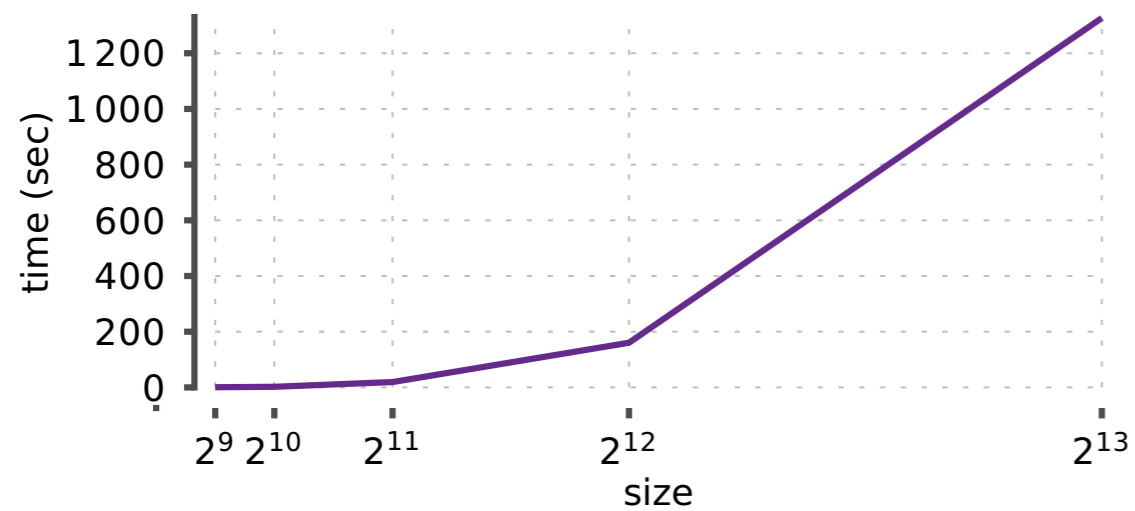


Example

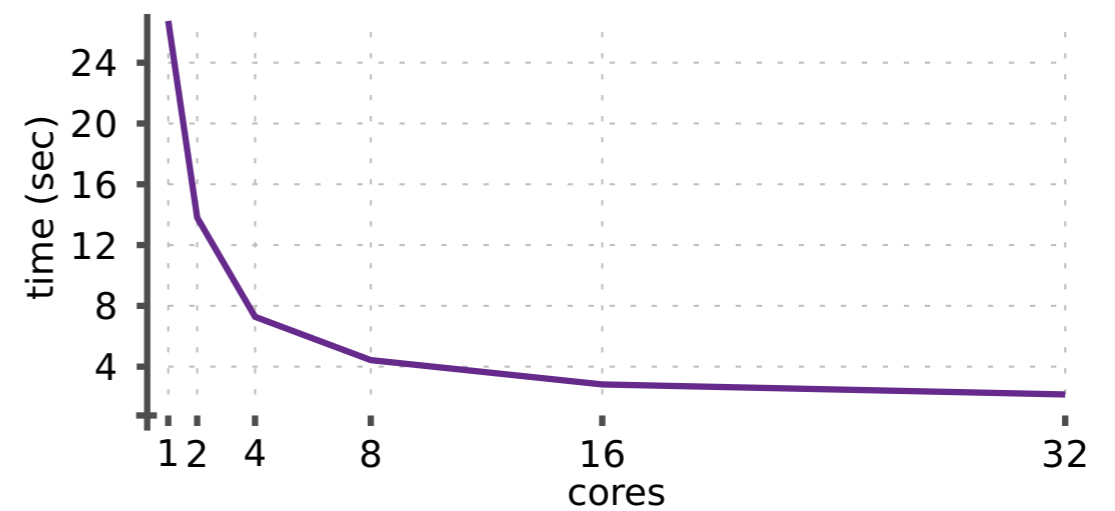
$$\begin{array}{c} 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{array} \begin{array}{c} a \\ b \end{array} \begin{array}{c} 1 \\ 0 \\ 1 \\ 1 \\ 0 \\ 0 \end{array}$$



Scalability



Scales well with input size



Parallelises well to many cores

Summary

- C1P and cyclic factors further improve the interpretability of BMF factors
 - Natural for edge bundle plots
- PQ-trees facilitate the efficient testing of compatible factors
- Possible extensions to other algebras, time series, ...

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Thank You!